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Contents

[Introduction 3](#_Toc58004131)

[What the Project is About 3](#_Toc58004132)

[Background Information on Socket Programming 3](#_Toc58004133)

[Description of the Client and Server Programs 3](#_Toc58004134)

[Basic Approach to Implement the Protocol 3](#_Toc58004135)

[Detailed Description of the Index Server Program 3](#_Toc58004136)

[Detailed Description of the Client Program 3](#_Toc58004137)

[Observations and Analysis 3](#_Toc58004138)

[Test 1: List Local Files Versus Display Online Registry 3](#_Toc58004139)

[Test 2: Have All Files Registered and Display Registry 3](#_Toc58004140)

[Test 3: Try to Register a File Already Registered to Same Username 3](#_Toc58004141)

[Test 4: Attempt to Download Files with Duplicated and Unique Username 3](#_Toc58004142)

[Test 5: De-Register Files with Original Username 3](#_Toc58004143)

[Test 6: Quitting the Program 3](#_Toc58004144)

[Conclusions 3](#_Toc58004145)

[References 3](#_Toc58004146)

[APPENDICES 3](#_Toc58004147)

[P2Pserver Source Code 3](#_Toc58004148)

[P2Ppeer Source Code 3](#_Toc58004149)

# Introduction

## What the Project is About

The objective of this project is to develop and implement the code for a basic peer-to-peer (P2P) file sharing application. The basic idea of a P2P file-sharing network is that many computers come together and pool their resources to form a content distribution system. The computers are called peers because each one can alternatively act as a client to another peer, fetching its content, and as a server, providing content to other peers [1]. P2P technology was used by popular services such as Napster and LimeWire. The most popular protocol for P2P sharing is BitTorrent [2].

For this project, a client host (i.e. representing a user) will query a server that provides an index of file names/addresses. If the file information is stored with the server, the server will send the information to the client. The client will then request the file from the content server that has the actual file. If the file download is successful, the client host will register its own address with the index server as another content server that could provide the file. Any content server is potentially a client and any client could potentially become a content server. The non-index server hosts are thus network peers.

The program will be implemented in C programming language because of its relative ease of implementation for network programming. It is simple with an easily readable syntax. It is very efficient, requiring fewer lines of code to complete more work.

## Background Information on Socket Programming

The concept of the socket is fundamental to network programming. Processes that want to communicate with one another need an interface. Transmission is governed by factors such as the address domain (e.g. IPv4), whether the data is streamed and requires a formal connection or segmented and connectionless, and the transport protocol. The specific IP address and port number of the receiving host must be indicated. All that information is accounted for by the socket application programming interface (API).

Analogous to its electrical hardware namesake (i.e. female connector), a network socket is a communication endpoint. It controls and organizes the communication traffic that a node could receive. Structurally it is a file descriptor and returns an integer, but it has instructions for interfacing with a node rather than working with a file. It stores information with respect to the protocol family (e.g. IPv4), the form of the sent data (continuous or segmented), and the governing transport protocol.

# Description of the Client and Server Programs

## Basic Approach to Implement the Protocol

To implement the protocol, the index server should be active first, to be available when a client wants to interact with it. Next, a peer with a file that other peers may want to download should be activated and have its file registered with the index server. After the file is registered it becomes accessible for download. Client peers access the file by querying the index server for the list of online registered files. The client user is then prompted for which file to download. If the file is registered with the index server, the index server arranges for the content server to send the file to the client, provided the file is available. The client peer then registers its possession of the file with the index server and thus becomes a content server itself. Future client peers searching for the file would be directed to the latest peer to register the file. When a content server wants to de-register a file, it would do so by contacting the index server.

Essentially, packets of information, protocol data units (PDUs), are being passed back and forth between hosts. Each PDU consists of a 100-byte payload portion, where data is stored, and a one-byte label portion, where the type of data being transmitted is identified by a letter. The payload portion may in turn have multiple subcomponents for multiple chunks of information. In this project, the PDU will be represented by a typedef struct data structure.

To transmit and receive PDUs, socket objects must be created and bound to the relevant host address. For secure, reliable transmission (i.e. TCP protocol), a formal connection must be made between the hosts via their sockets. For insecure/less reliable transmission (i.e. UDP protocol), a formal connection is not made. Structures like typedef cannot be transmitted through a socket as a socket will only accept binary data. The PDUs therefore must be serialized into byte form for transmission, then de-serialized back into object form by the receiver for interpretation. The PDUs were sent by invoking ‘&rpdu’.

When a PDU is received, the receiving host must make decisions based largely on they label of the PDU and/or the presence or not of data. A table of instructions based on the PDU label is given below.

**Table 1: Client Directions based on PDU Label**

|  |  |  |
| --- | --- | --- |
| **PDU Type** | **Function** | **Direction** |
| R | Content Registration | Peer to Index Server |
| D | Content Download Request | Content Client to Content Server |
| S | Search for content and the associated content server | Between Peer and Index Server |
| T | Content De-Registration | Peer to Index Server |
| C | Content Data | Content Server to Content Client |
| O | List of On-Line Registered Content | Between Peer and Index Server |
| A | Acknowledgement | Index Server to Peer |
| E | Error | Between Peers or between Peer and Index Server |

Received bytes are read or written (as applicable) continuously in a while-loop. Once there is no more data, or in some cases when a timeout exception is triggered, the reading/writing operation stops. If no data is received at all when it is expected, an exception is triggered, and an error message is returned to the sender. Communication between a peer and the index server is via the UDP protocol. It is connectionless and provides no guarantee of successful or error-free transmission, but it is faster. The trade-off is considered acceptable given that data between the server and peers consist mainly of messages, and not files. Transmission between peers involves file transfer, so the more reliable TCP protocol, which is connection-oriented and checks for errors, was used for peer-to-peer traffic.

Various specialized functions were called to implement the protocol. To access those functions, several external modules/libraries had to be imported. The socket module was discussed above. The table[MAXCON] structs was made to keep track of the registered content. The function memset() was used to fill a block of memory with a particular value and to setup a UDP Connection with the index server. The socket() function was used to represent a file descriptor and allocate a socket for the index server. The Initialization of select structure and table structure was done through: the use of FD\_zero() was made to initialize the file descriptor; FD\_SET() was used for listening on the index server socket and listening on the read descriptor.

## Detailed Description of the Index Server Program

The index server is responsible for maintaining the list of online registration content, managing the registering and de-registering of files, and directing search requests to the appropriate content server.

The index server code starts with preliminary definitions such as defining the PDU and the files list as types of typedef structs, declaring an empty (initially) array for the server file registration list, setting the port number and host ID, and creating a UDP socket object. A key functionality of the server in this project is its ability to handle multiple client requests/processes simultaneously.

The index server’s operation begins by binding the socket object to the host’s IP address and port number. If an error occurs, an exception is triggered. After binding the socket, the server indicates that it is waiting for a connection and starts listening for potential clients. If a connection request is accepted and a connection is made, the server will print out the IP address and port number of the host that it is connected to.

Once connection is established. The packet is received in binary form then de-serialized back into a PDU typedef struct. A message is printed acknowledging receipt of the specific PDU from a peer. The label of the PDU is checked for PDU type. After binding the socket, The type could be “Content Registration Request- R”, “Search Content- S”, “List current Content- O” or “De-registration- T”.

If the type is “R”, then the peer name, filename, and peer address are extracted and stored in variables. A “registration()” function is called with parameters ‘s, rpdu.data, & fsin’ (to assume that the file to be registered does not already exist in the registry). Every index in the server file registration list is matched against the connected peer name–file name combination. If there is a match(list[n].head!=NULL), then the “find” variable is set to “n” and a PDU with an error message is sent back to the client telling them to pick another username to register that file. If the file does not exist, then the peer name, filename, and peer address are added to the server file registration list and an acknowledgement PDU is sent back to the client.

If the type is “S”, then the file name is extracted and stored in a variable. A “search\_content()” function is called with parameters ‘s, rpdu.data, & fsin’. A for-loop is run to search the file name. If the result is a negative integer then a message is print “Search Unsuccessful” and tpdu.type is set to 'E'. If there is a match, then a while loop is run and the tpdu.type is set to 'S' and the address of the matched file name is stored. An S-type acknowledgement PDU is sent to the client and a message stating that the PDU is being sent is printed. If the content is not found, then a PDU with an error message is sent to the client.

If the type is “O”, then the server first checks to see if there are any entries in the server file registration list. If there are no entries, then an O-type PDU is returned to the client. A message stating that such a PDU is being sent to the peer is printed. If there are entries in the list, then the peer name, filename, and peer address of every entry is sent to the client in an O-type PDU.

If the type is “T”, then the peer name, filename, and peer address are also extracted and stored in variables. A “deregistration()” function is called with parameters ‘s, rpdu.data, & fsin’. Every index in the server file registration list is matched against the connected peer name–file name combination. If there is a match, then the matching entry is removed from the server file registration list, and acknowledgement PDU is sent to the client, and a message stating that the file has been deregistered to remove the entries is printed out. When there are no more entries in the server file registration list (i.e. the length of the list equals loop count), then an error PDU is sent to the client stating that the file does not exist.

## Detailed Description of the Client Program

The P2P client sends connection requests to the index server. It is responsible for sending requests to the index server to register or de-register a file. It is responsible for requesting the list of online registered content files and choosing which file it wants to download. It is responsible for requesting the list of local registered content files. Finally, it is responsible for quitting an online session (after deregistering its files). The P2P client may also function as a content server. In order to accommodate concurrent operation as client and server.

Like the server code, the P2P client code also starts with preliminary definitions and declarations. A UDP socket is created. The port number and host ID are set. A socket timeout is set. The socket is connected to the index server host ID and port number. Finally, a PDU is created as a typedef structs (with the same structure as in the code for the index server).

The code for the P2P client is structured with function definitions: search\_content(), deregistration(), registration(), client\_download(), server\_download(). The actual operation begins with a prompt for the user to input a username. The client then starts both client and server function and executes its commands.

The function begins by prompting the user to choose a client function from the list: “O” (get online list); “L” (list local files); “R” (register a file); “T” (de-register a file); or “Q” (quit the program). The search\_content() function takes the following parameters: int s\_sock, char \*name, PDU \*rpdu. It sets the tpdu.type to 'S' and copies it onto the memory using the memcpy() function. Apart from searching the content and reading and writing error reports, it will print “Search\_Content: Cannot find content” for rpdu->type == 'E', “Search: Find content” for rpdu->type == 'S', and “Search\_Content: Protocol Error” if neither of these cases are true.

If the user selects “O”, then an “online\_list()” function is called with the parameter s\_sock(socket) and an O-type PDU is sent to the index server. The responding PDU from the index server is downloaded and de-serialized until a timeout exception is triggered, signalling download completion. If the received PDU is Type-O, then its payload is checked. If there are entries in the list, then the peer name and file name of each entry is displayed, and the user is prompted for which file to download. A message is printed “On-line Content List:”.

An S-type PDU with the filename request as the payload is then created and sent to the index server. As before, the server’s response is downloaded and serialized until a timeout exception is triggered. If the responding PDU is also Type-S, then the address of the peer where the file is stored is extracted from the payload and the server\_download() function is called, with the peer address, filename request, and file index list as arguments.

If the user selects “L”, then the local\_list() function is called with the folder containing the files as an argument, and the files are printed out with the message “Local Content List”. This function searches up the table of entries and upon a match, displays the file.

If the user selects “R”, then the user is prompted for the name of the file to register via a message “Enter the name of the content” and then the registration() function is called with the username, filename index request, and peer port number as arguments.

If the user selects “T”, then the user is prompted for the name of the file to de-register through a message “Enter the name of the Content”. If the file exists, the deregistration() function is called with the same arguments as the registration() function, except the filename index request parameter is replaced with a filename index delete request parameter. If the file does not exist or is not associated with the user, a message stating such is printed out.

A P2P client can also function as a content server if it has registered a file with the index server. Preliminary declarations include the creation of a TCP socket for communication between peers, a socket timeout set to one second, and empty arrays for outputs, exceptions, and the sockets of peers requesting a connection.

The content server will listen for available sockets which are ready through the FD\_SET() function. The content server will check the socket list for sockets to establish a connection. When a new connection is made with a peer, the peer gets added to the socket list.

If it is Type-D, the content server will check the directory for the file. If the file is there, it will be opened, read, and sent to the client peer that requested it in a PDU of Type-C. The file would then be closed and a notification stating that transmission is complete, and the connection will be closed, will be printed.

If the content server cannot find the file, an error PDU (i.e. E-type) with a “File does not exist” message would be sent back to the client peer and the connection would be closed.

If the item in the socket list is not incoming data, it would be removed from the socket list.

The P2P client implements three functions: file registration (registeration()), file download (server\_download()), and file de-registration (deregisteration()). The file registration function takes two parameters: filename and peer port. When the function is called, an R-type PDU is created and sent to the index server. The response packet from the index server is downloaded and de-serialized until there is no more data. The PDU type is then checked. If it is Type-A (an acknowledgement PDU), the filename is added to the file index and the payload’s contents is printed out. If the response was a Type-E PDU, the error message is printed out and the user is prompted for a different username to send with a new file registration request. The new R-type PDU is prepared and sent.

# Observations and Analysis

The following are tests of basic functionality.

## Test 1: List Local Files Versus Display Online Registry

This is a reasonable preliminary test to show the difference between a file existing and a file being registered. The first command is expected to succeed. The second command is expected to fail.

***Figure 1: List of local files versus online registry when no files have registered.***

The test succeeded. The local files exist and can be displayed, but there are no files registered so the online registry cannot be displayed.

## Test 2: Have All Files Registered and Display Registry

With entries in the registry, a call to view the registry or download files should be successful. The test is expected to succeed.

***Figure 2: Online registry with all files registered.***

The test was successful.

## Test 3: Try to Register a File Already Registered to Same Username

The program is designed to only allow one username to register a file. The attempt should fail, and an error message should be triggered.

***Figure 3: Registering a file already registered to the same username.***

The test was successful.

## Test 4: Attempt to Download Files with Duplicated and Unique Username

The code was designed to block peers with the same username from downloading the same file but allow peers with a different username to download and register the same file. The test should be successful.

***Figure 4: Downloading of file with duplicated username and unique username.***

The test was successful. A second user named Leonardo had to change his name (to Dante) to download file1.txt, which was already registered to the first Leonardo. After registering file1.txt under the changed name, the other files were registered under the original name, with no user-input reversion to the original name.

## Test 5: De-Register Files with Original Username

The second Leonardo should be able to de-register his files. However, his registration of file1.txt was under the name “Dante”. This test will see whether the user at port 6050 (Leonardo/Dante) can still access his registration under the temporary name “Dante”.

***Figure 5a: Deregistration of files under original username.***

De-registration of the files under the name “Leonardo” was successful.

***Figure 5b: Deregistration of file registered under temporary (and inaccessible) username.***

## Test 6: Quitting the Program

Quitting the program automatically executes the de-register file(s) function beforehand.

***Figure 6a: Demonstration of the quit() functionality.***

In this scenario, all users except the second Leonardo username have quit. The files of other users are no longer registered after they quit the program. The second Leonardo has been frozen because he tried to de-register a file listed under a temporary name. Trying to de-register a file under a username that is not the one the file is registered would normally trigger an error stating that only usernames registered to a file could de-register that file. The error in this case is different.

***Figure 6b: Demonstration that files registered under a temp name are inaccessible by the present code.***

# Conclusions

The specs were largely met, but not without challenges or room for improvement. The objective of this project was to develop and implement a peer-to-peer file-sharing application. The basic functionalities of client-to-index server interaction and peer-to-peer interaction were achieved. A client can register and de-register files with the server. A client can request to see an online registry and download registered files from a peer (“content server”). clients can request to see the local file directory. Peers can quit the application and have their registered files automatically downloaded. A peer can run client and server processes simultaneously, a functionality that was particularly challenging to incorporate.

Despite the generally positive results, the tests also revealed a setback. If a client tried to register or download a file that was already registered to a peer with the same username, the client was supposed to be blocked and prompted to change their username. However, a username change was not permanent. Files without a name conflict that were subsequently registered by the peer were done so under the original username. The system no longer recognized the substitute username as belonging to the peer, so the file registered under that substitute username could not be accessed to be deregistered. By contrast, the files registered under the original username could be deregistered. In summary, the system has implementation problems because it was designed to manage files based on username, but it was not designed to keep consistent track of mandatory username changes.

# References

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# APPENDICES

## P2Pserver Source Code

## P2Ppeer Source Code